Row-factorization matrices and type of almost Gorenstein monomial curves

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Preliminaries

Notation:

- ${\mathcal S}$ is a numerical semigroup (submonoid of ${\mathbb N}$ such that ${\mathbb N}\setminus {\mathcal S}$ is finite)
- $\mathcal{G} = \{g_1, \dots, g_e\}$ are the minimal generators \mathcal{S} ;
- $e = |\mathcal{G}|$ is the embedding dimension of \mathcal{S} ;
- $F(S) = \max \mathbb{Z} \setminus S$ is the Frobenius number of S;
- The set of pseudo-Frobenius numbers

 $PF(\mathcal{S}) = \{ x \notin \mathcal{S} \mid x + g_i \in \mathcal{S} \text{ for all } i = 1, \dots, e \}.$

• t(S) = |PF(S)| is the type of S.

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Main topic

Study the boundedness of t(S) in terms of e.

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In the general case of semigroup rings it is known that:

- **1** if e = 2, then t(S) = 1.
- 2 if e = 3, then $t(S) \leq 2$.

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If e = 4

Example (Bresinsky, 1975)

The family of numerical semigroups

$$\mathcal{S}_h = \langle (2h-1)2h, (2h-1)(2h+1), 2h(2h+1), 2h(2h+1) + 2h - 1 \rangle$$

is such that $t(\mathcal{S}_h) = 4h - 3$.

 \implies In the general case the type is not bounded.

We say that S is **almost symmetric** if, for every $x \notin S$, we have either $F(S) - x \in S$ or $\{x, F(S) - x\} \subseteq PF(S)$.

Question (Numata, 2013)

Let S be an almost symmetric numerical semigroup with e(S) = 4. Is it true that $t(S) \le 3$?

Families of almost symmetric semigroups such that t(S) is large are present in the literature [García-Sánchez and Ojeda, 2019], but for all of them we have $2e \ge t$. We say that a matrix $\Lambda = (\lambda_{ij}) \in M_e(\mathbb{Z})$ is a **RF-matrix** for $f \in PF(S)$ if for every i = 1, ..., e, $\lambda_{ii} = -1$ $\lambda_{ij} \in \mathbb{N}$ if $i \neq j$ and $\lambda_{i1}g_1 + ... + \lambda_{ie}g_e = f$.

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Example

Take $S = \langle 6, 7, 9, 10 \rangle$, $PF(S) = \{3, 8, 11\}$. The two matrices $\Lambda_1 = \begin{pmatrix} -1 & 2 & 0 & 0 \\ 1 & -1 & 1 & 0 \\ 0 & 1 & -1 & 1 \\ 0 & 0 & 2 & -1 \end{pmatrix}, \quad \Lambda_2 = \begin{pmatrix} -1 & 2 & 0 & 0 \\ 1 & -1 & 1 & 0 \\ 0 & 1 & -1 & 1 \\ 3 & 0 & 0 & -1 \end{pmatrix}$ are both RF-matrices for $8 \in PF(S)$.

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Proposition

Let $\Lambda_1 = (a_{ij})$ be a RF-matrix for f and $\Lambda_2 = (b_{ij})$ a RF-matrix for F(S) - f. Then for every $i \neq j$ we have $a_{ij}b_{ji} = 0$.

Example

Take $S = \langle 6, 7, 9, 10 \rangle$. We have $PF(S) = \{3, 8, 11\}$. The matrix Λ_1 is a RF-matrix for 3, while Λ_2 is a RF-matrix for 8 = F(S) - 3.

$$\Lambda_1 = egin{pmatrix} -1 & 0 & 1 & 0 \ 0 & -1 & 0 & 1 \ 2 & 0 & -1 & 0 \ 1 & 1 & 0 & -1 \end{pmatrix}, \quad \Lambda_2 = egin{pmatrix} -1 & 2 & 0 & 0 \ 1 & -1 & 1 & 0 \ 0 & 1 & -1 & 1 \ 0 & 0 & 2 & -1 \end{pmatrix}.$$

Simple case: e = 4

Let
$$i, j \in \{1, \dots, e\}$$
, $i \neq j$ and
 $m_{ij} = max\{K \in \mathbb{N} \mid Kg_j - g_i \notin S\}$, $M_{ij} = m_{ij}g_j - g_i \notin S$,
and $\mathcal{M} = \{M_{ij} | i \neq j\}$, $|\mathcal{M}| \leq e(e-1)$.

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and $\mathcal{M} = \{M_{ij} | i \neq j\}, |\mathcal{M}| \leq e(e-1).$
Let $f, F(S) - f \in PF(S).$
• If $f = kg_i - g_i \in PF(S)$, then $f \in \mathcal{M}$.

- If there is a row with at least e 2 zeroes, then the element associated to that RF-matrix belongs to M.
- In every couple of RF-matrices there are at least e(e 1) zeroes distributed over 2e rows.

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Theorem

If S is almost symmetric and e = 4 then t(S) is bounded.

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Theorem

If S is almost symmetric and e = 4 then t(S) is bounded. (Actually, $t(S) \le 3$)

Computational evidence [Garcia-Sanchez]

If S is almost symmetric, e = 5 and $g_5 \leq 200$, then $t(S) \leq 5$.

If e = 5 then e(e - 1) = 20 = 2e(e - 3), so the previous argument leaves out one case:

• Every row and column of Λ_1, Λ_2 has exactly 2 zeroes.

Example

For instance we could have

$$\Lambda_1 = \begin{pmatrix} -1 & 0 & 0 & * & * \\ 0 & -1 & 0 & * & * \\ * & * & -1 & 0 & 0 \\ 0 & * & * & -1 & 0 \\ * & 0 & * & 0 & -1 \end{pmatrix} \ \Lambda_2 = \begin{pmatrix} -1 & * & 0 & * & 0 \\ * & -1 & 0 & 0 & * \\ * & * & -1 & 0 & 0 \\ 0 & 0 & * & -1 & * \\ 0 & 0 & * & * & -1 \end{pmatrix}$$

Lemma

For every possible distribution of zeroes of the form described before, there are at most two elements of PF(S) having a RF-matrix of that shape.

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If S is almost symmetric and e = 5 then t(S) is bounded. (Actually, $t(S) \le 473$)

• If e = 4 we could easily determine the shape of the factorization of all elements in PF(S), showing that $f = kg_i - g_j$.

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- If e = 5 we could bound the number of elements of PF(S) associated to a fixed shape of the RF-matrix.
- If $e \ge 6$ all of this does not work anymore.
 - There are a lot of elements not of the form $f = kg_i g_j$.
 - It seems that there are more elements associated to RF-matrices with the same shape.
- It is not clear whether t(S) is bounded if $e \ge 6$.

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Example

Take $S = \langle 455, 497, 574, 589, 631, 708 \rangle$. Then t(S) = 14 and

 $PF(S) = \{3079, 3289, 3521, 3655, 3674, 3789, 3923, 4057, \}$

4172, 4191, 4325, 4557, 4767, 7846}.

The elements $\{3521, 3655, 3789, 3923, 4057, 4191\} \subset PF(S)$ all have RF-matrices sharing the same shape.

Some remarks:

- The example above is the first known example of almost symmetric numerical semigroup such that t > 2e.
- "Bad" examples occur for very high values of g_i very hard to find by computer.

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The previous example was built using this construction. It is not known if this can be done for arbitrarily large sets of RF-matrices.

Thank you for your attention.