

# **Optimierung** I

## Ubungsblatt 3 Bearbeitung bis 1. April 2014

#### Aufgabe 3.1: [Canonical form]

Let  $A \in \mathbb{R}^{m \times n}$ ,  $b \in \mathbb{R}^m$ ,  $c \in \mathbb{R}^n$ . Put the following linear problem into a canonical form:

$$\min_{x \in \mathbb{R}^n} c^\top x, \text{ subject to } Ax = b, \ x \ge 0, \ x_1 \le x_2 \le \dots \le x_n.$$
(1)

## Aufgabe 3.2: [Piecewise affine approximation of a convex function]

Let  $A \in \mathbb{R}^{m \times n}$ ,  $b \in \mathbb{R}^m$ ,  $c \in \mathbb{R}^n$ . Let  $f : \mathbb{R}^n \to \mathbb{R}$  be a differentiable convex function. We are interested in the following non-linear optimization problem:

$$\min_{x \in \mathbb{R}^n} f(x), \text{ subject to: } Ax = b, \ x \ge 0.$$
 (P)

We assume that the feasible set is bounded. Let s be the optimal value of problem (P). We recall that since f is convex, then for all x and x' in  $\mathbb{R}^n$ ,

$$f(x') - f(x) \ge Df(x)(x' - x),$$
 (2)

where Df(x) is the derivative of f. Let  $x_1,...,x_N$  be N points of  $\mathbb{R}^n$ . We define:

$$f_N(x) = \max_{i=1,\dots,N} \{ f(x_i) + Df(x_i)(x - x_i) \}$$
(3)

and we consider the problem

$$\min_{x \in \mathbb{R}^n} f_N(x), \text{ subject to: } Ax = b, \ x \ge 0.$$
 (P<sub>N</sub>)

We consider that  $f_N$  is a piecewise affine approximation of f and that problem  $(P_N)$  is an approximation of (P).

- i) For a function f and points  $x_1, \dots, x_N$  of your choice, draw f and  $f_N$ .
- ii) Prove that for all  $x \in \mathbb{R}^n$ ,  $f(x) \ge f_N(x)$ .
- iii) Compute  $f_N(x_i)$ , for all *i*.
- iv) Formulate problem  $(P_N)$  as a linear programming problem, justify the reformulation. Put the obtained problem into a canonical form.
- v) Let  $x^*$  be a solution to problem  $(P_N)$ , prove that

$$f_N(x^*) \le s \le f(x^*). \tag{4}$$

vi) How could we iteratively improve the approximation  $(P_N)$  of problem (P)? Describe an algorithm which enables to solve approximately problem (P) by solving only linear programming problems.

### Aufgabe 3.3: [Application exercise]

A manufacturer wishes to product a special alloy (Legierung), made of a proportion  $b_1$  of a metal 1 and a proportion  $b_2$  of a metal 2 (with  $b_1 + b_2 = 1$ ). To this purpose, he can mix n different alloys. The *i*-th alloy is made of a proportion  $a_{1,i}$  and  $a_{2,i}$  of metals 1 and 2 respectively, with  $a_{1,i} + a_{2,i} = 1$ . The prices (per unit of weight) of the n alloys are  $c_1, c_2, ..., c_n$ .

- i) Denoting by  $x_i$  the proportion of alloy *i* used to make the special alloy, formulate the manufacturer's problem as a linear programming problem. Warning: use as few constraints as possible.
- ii) What is the maximal number of bases of the problem ?
- iii) Give a condition on the data to ensure that the problem has a solution.
- iv) We assume that  $a_{1,1} \leq a_{1,2} \leq \dots \leq a_{1,n}$ . We assume that for some  $k, a_{1,k} < b_1 < a_{1,k+1}$ . What is the number of feasible bases ?

	Alloy 1	Alloy 2	Alloy 3	Alloy 4	Special alloy
Metal 1	0.2	0.3	0.5	0.7	0.4
Metal 2	0.8	0.7	0.5	0.3	0.6
Cost per unit	5	3	7	6	

v) We consider a case with 4 alloys:

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## Aufgabe 3.4: [Optimal transportation]

Quantities  $a_1,...,a_m$ , respectively, of a certain product are to be shipped from each of m locations and received in amounts  $b_1, b_2,...,b_n$ , respectively at each of n destinations. Associated with the shipping of a unit of product from origin i to destination j is a unit shipping cost  $c_{ij}$ . It is desired to determine the amounts  $x_{i,j}$  to be shipped between each origin-destination pair i = 1,...,m, j = 1,...,n so as to satisfy the shipping requirements and minimize the total cost of transportation. To make the problem consistent, we assume that  $\sum_{i=1}^m a_i = \sum_{j=1}^n b_j$ .

i) Formulate the problem as a linear programming problem. Denoting by

$$y = (x_{1,1}, \dots, x_{1,n}, x_{2,1}, \dots, x_{2,n}, \dots, x_{m,1}, \dots, x_{m,n})$$
(5)

the vector of size nm, write the constraints of the problem in the canonical form Ay = w, with  $w \in \mathbb{R}^{(n+m)\times 1}$  and  $A \in \mathbb{R}^{(n+m)\times nm}$ .

- ii) Is it possible that the rank of the matrix A is equal to n + m?
- iii) We consider the following case: m = 2, n = 3, a = (5, 8), b = (4, 3, 6), with the costs:

	Destination 1	Destination 2	Destination 3
Origin 1	1	1	2
Origin 2	2	1	1

Express the variables  $x_{1,2}$ ,  $x_{1,3}$ ,  $x_{2,1}$ , and  $x_{2,3}$  in function of  $x_{1,1}$  and  $x_{2,2}$ . Reduce the optimal transportation problem into a problem with two optimization variables. Solve it graphically.