

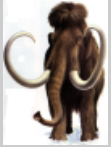
Musical Enumeration Theory Applied to the Classification of Canons

Harald Friepertinger

Karl-Franzens-Universität Graz

Colloquium on Mathematical Music Theory

University Graz, May 6 – May 9, 2004



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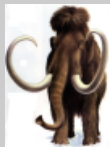
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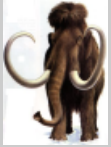
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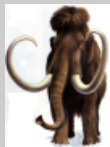
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- basic facts about classification of discrete structures, especially introducing group actions,



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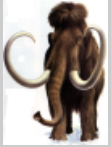
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- examples of discrete structures in music theory,



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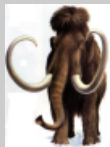
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- what are discrete structures,
- basic facts about classification of discrete structures, especially introducing group actions,
- examples of discrete structures in music theory,
- some results on the enumeration of canons.



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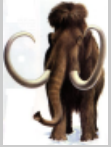
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Discrete Structures



Discrete structures are objects which can be constructed as:

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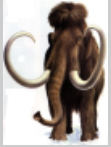
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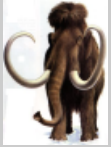
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Discrete Structures

Discrete structures are objects which can be constructed as:
– subsets, unions, products of finite sets,



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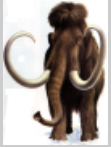
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Discrete Structures

Discrete structures are objects which can be constructed as:

- subsets, unions, products of finite sets,
- mappings between finite sets,



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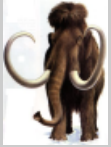
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Discrete Structures

Discrete structures are objects which can be constructed as:

- subsets, unions, products of finite sets,
- mappings between finite sets,
- bijections, linear orders on finite sets,



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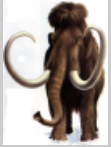
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Discrete Structures

Discrete structures are objects which can be constructed as:

- subsets, unions, products of finite sets,
- mappings between finite sets,
- bijections, linear orders on finite sets,
- equivalence classes on finite sets,



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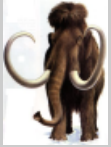
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Discrete Structures

Discrete structures are objects which can be constructed as:

- subsets, unions, products of finite sets,
- mappings between finite sets,
- bijections, linear orders on finite sets,
- equivalence classes on finite sets,
- vector spaces over finite fields, . . .



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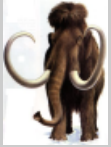
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- vector spaces over finite fields, . . .

Examples: graphs, necklaces, designs, codes, matroids, switching functions, molecules in chemistry, spin-configurations in physics, objects of local music theory.

Classification of Discrete Structures I



The process of ***classification*** provides more detailed information about the objects in a discrete structure. We distinguish different steps in the process of classification:

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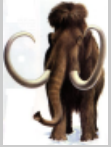
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Classification of Discrete Structures I



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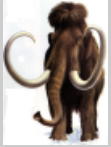
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The process of ***classification*** provides more detailed information about the objects in a discrete structure. We distinguish different steps in the process of classification:

step 1: Determine the number of different objects.

Classification of Discrete Structures I



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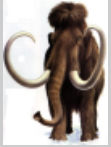
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The process of ***classification*** provides more detailed information about the objects in a discrete structure. We distinguish different steps in the process of classification:

step 1: Determine the number of different objects.

step 2: Determine the number of objects with certain properties.

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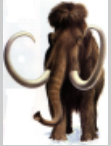
The process of **classification** provides more detailed information about the objects in a discrete structure. We distinguish different steps in the process of classification:

step 1: Determine the number of different objects.

step 2: Determine the number of objects with certain properties.

step 3: Determine a complete list of the elements of a discrete structure.

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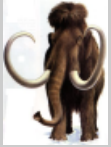
The process of **classification** provides more detailed information about the objects in a discrete structure. We distinguish different steps in the process of classification:

step 1: Determine the number of different objects.

step 2: Determine the number of objects with certain properties.

step 3: Determine a complete list of the elements of a discrete structure.

step 4: Generate the objects of a discrete structure uniformly at random.



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Labelled Graphs

A simple labelled graph is described by a finite set of labelled vertices V and by a set of edges connecting two vertices. An edge is usually described as a 2-set of vertices, indicating which two vertices are connected by the corresponding edge.



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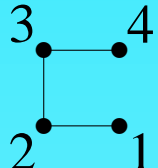
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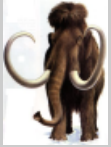
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For instance $V = \{1, 2, 3, 4\}$ and $E = \{\{1, 2\}, \{2, 3\}, \{3, 4\}\}$ describes

the graph  on 4 vertices.



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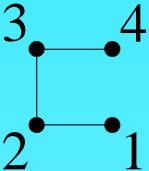
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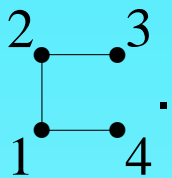
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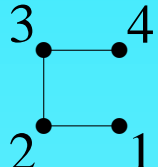




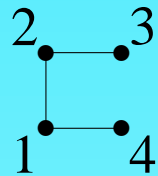
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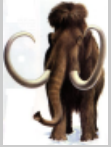
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important, whence this graph is considered to be essentially the same as

. All graphs obtained by relabelling of one graph are collected to an unlabelled graph, an equivalence class of labelled graphs. In our example

we get .



Unlabelled Graphs on 4 Vertices

As we have just seen, often the elements of a discrete structure are themselves classes of objects which are considered to be equivalent. In one class all those elements are collected which are assumed to be not essentially different. These classes are caused by relabelling of labelled structures or by otherwise naturally motivated equivalence relations.

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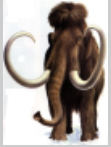
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Example: Classification of unlabelled simple graphs:

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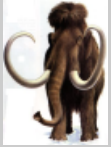
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Example: Classification of unlabelled simple graphs:
step 1: There are 11 graphs on 4 vertices.

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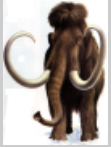
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Unlabelled Graphs on 4 Vertices

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Example: Classification of unlabelled simple graphs:

step 1: There are 11 graphs on 4 vertices.

step 2: There exists exactly one graph with 0, 1, 5 or 6 edges; two graphs with 2 or 4 edges; three graphs with 3 edges.

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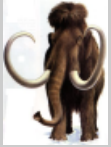
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Unlabelled Graphs on 4 Vertices

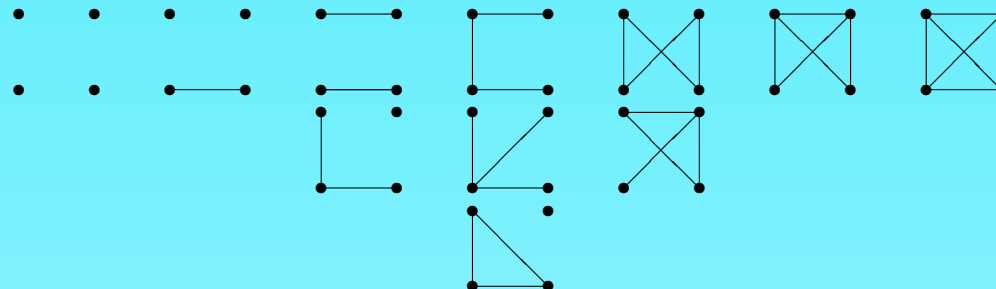
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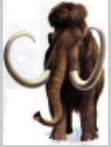
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step 3:



Classification of Discrete Structures II



1. Determine the finite set X on which the discrete structure is defined.

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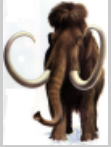
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Classification of Discrete Structures II



1. Determine the finite set X on which the discrete structure is defined.
2. Describe how the discrete structure is constructed over X .

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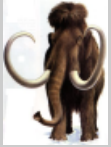
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Classification of Discrete Structures II



1. Determine the finite set X on which the discrete structure is defined.
2. Describe how the discrete structure is constructed over X .
3. If the objects of a discrete structure are equivalence classes, introduce a suitable group action on X to describe the equivalence classes as orbits.

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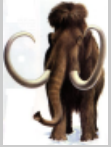
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Classification of Discrete Structures II



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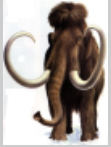
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1. Determine the finite set X on which the discrete structure is defined.
2. Describe how the discrete structure is constructed over X .
3. If the objects of a discrete structure are equivalence classes, introduce a suitable group action on X to describe the equivalence classes as orbits.
4. Proceed with the four steps for classification (enumeration, construction, random generation) as described above.



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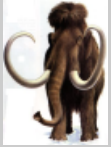
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Group Actions

A multiplicative group G with neutral element 1 acts on a set X if there exists a mapping

$$*: G \times X \rightarrow X \quad * (g, x) \mapsto g * x$$



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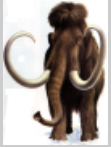
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Notation: We usually write gx instead of $g * x$, and $\bar{g}: X \rightarrow X$, $\bar{g}(x) = gx$.

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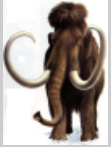
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Notation: We usually write gx instead of $g * x$, and $\bar{g}: X \rightarrow X$, $\bar{g}(x) = gx$.

A group action will be indicated as ${}_G X$.

If G and X are finite sets, then we speak of a **finite group action**.

Group Actions



Orbits under Group Actions

A group action $G \curvearrowright X$ defines the following equivalence relation on X .
 $x_1 \sim x_2$ if and only if there is some $g \in G$ such that $x_2 = gx_1$.

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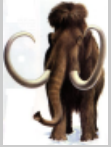
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Orbits under Group Actions

A group action $G \curvearrowright X$ defines the following equivalence relation on X .
 $x_1 \sim x_2$ if and only if there is some $g \in G$ such that $x_2 = gx_1$. The
equivalence classes $G(x)$ with respect to \sim are the **orbits** of G on X .
Hence, the orbit of x under the action of G is

$$G(x) = \{gx \mid g \in G\}.$$

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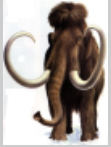
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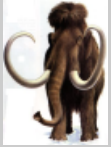
Orbits under Group Actions

A group action $G \curvearrowright X$ defines the following equivalence relation on X .
 $x_1 \sim x_2$ if and only if there is some $g \in G$ such that $x_2 = gx_1$. The
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Hence, the orbit of x under the action of G is

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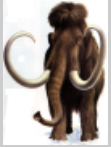
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Theorem. *The equivalence classes of any equivalence relation can be represented as orbits under a suitable group action.*



Stabilizers and Fixed Points

Let ${}_G X$ be a group action. For each $x \in X$ the **stabilizer** G_x of x is the set of all group elements which do not change x , in other words

$$G_x := \{g \in G \mid gx = x\}.$$

It is a subgroup of G .

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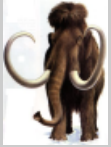
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Stabilizers and Fixed Points

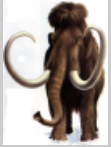
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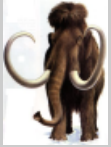
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Enumeration under Group Actions

Let $G \curvearrowright X$ be finite group action. The main tool for determining the number of different orbits is the

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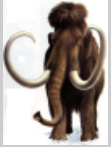
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Enumeration under Group Actions

Let $G X$ be finite group action. The main tool for determining the number of different orbits is the

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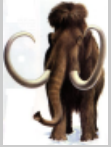
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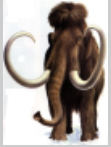
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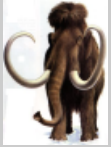
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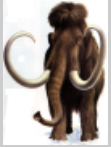
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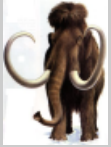
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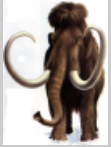
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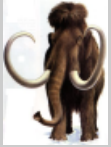
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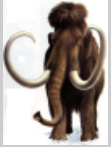
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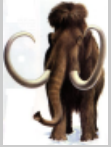
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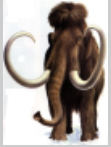
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The most important applications of classification under group actions can be described as mappings between two sets. Group actions on the domain X or range Y induce group actions on Y^X .

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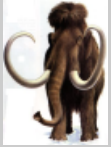
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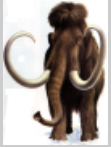


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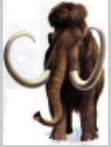
— Then H acts on Y^X by

$$H \times Y^X \rightarrow Y^X, \quad (h, f) \mapsto \bar{h} \circ f.$$

— Then the direct product $H \times G$ acts on Y^X by

$$(H \times G) \times Y^X \rightarrow Y^X, \quad ((h, g), f) \mapsto \bar{h} \circ f \circ \bar{g}^{-1}.$$

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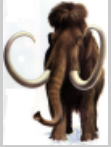
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Reduction of all pitches modulo one octave yields the notion of ***pitch-classes***. In an ***n-scale*** there are exactly n tones in one octave, whence there are n pitch classes which are described as elements of $\mathbb{Z}_n := \mathbb{Z}/n\mathbb{Z}$, the ***residue class ring modulo n*** .

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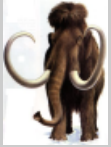
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Musical transposing or inversion can be described as ***transposing by one pitch-class, T*** , or ***inversion at pitch-class 0 , I*** , which are permutations on Z_n . They motivate cyclic groups, $\langle T \rangle$, dihedral groups, $\langle T, I \rangle$, or affine groups as permutation groups acting on Z_n .

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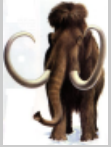
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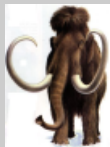
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Mosaics are partitions of Z_n . These are collections π of subsets of Z_n , such that the empty set is not an element of π and such that for each $i \in Z_n$ there is exactly one $P \in \pi$ with $i \in P$.

Discrete structures in Music Theory (cont.)



Next to the pitch component we also introduce the time as a second parameter.

For $n \geq 3$ a **tone-row** in Z_n is a bijective mapping $f: \{1, \dots, n\} \rightarrow Z_n$ where $f(i)$ is the pitch class of the tone which occurs in i -th position in the tone-row. Usually two tone-rows f_1, f_2 are considered to be equivalent if f_1 can be constructed by transposing, inversion and retrograde inversion R of f_2 . Thus the similarity classes of tone-rows are the $D_n \times \langle R \rangle$ orbits on the set of these bijections.

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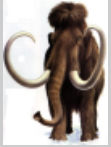
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In the most simple model of a motive, each note is described by its onset and by its pitch. In other words, we put a coordinate system over a line of notes. If there are exactly m possible onsets and n pitch classes, a **k -motive** is considered to be a k -subset of $Z_m \times Z_n$. In the case $m = n$, Mazzola showed that the group of all affine mappings on Z_n^2 is a musically motivated group.

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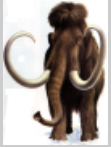
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Rhythms

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We also assume that we found a subdivision of the rhythm, i.e. a regular pulsation, into equidistant beats such that all rhythmical events coincide with some of these beats. If the rhythm is covered by a pulsation of n beats, then it can be described as a 0, 1-vector of length n , whence as a function $f: \{0, \dots, n-1\} \rightarrow \{0, 1\}$.

For example, the function $f = (f(0), \dots, f(7)) = (1, 0, 0, 1, 1, 0, 1, 0)$

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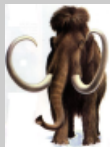
0,1-Vectors

The cyclic group C_n generated by $\pi_n := (0, 1, \dots, n-1)$ acts on the set of all mappings from $\underline{n} := \{0, 1, \dots, n-1\}$ to $\{0, 1\}$ according to

$$C_n \times \{0, 1\}^n \rightarrow \{0, 1\}^n, \quad (\sigma, f) \mapsto f \circ \sigma^{-1}.$$

If we write f as a vector $f = (f(0), \dots, f(n-1))$, then

$f \circ \pi_n^j = (f(n+j), \dots, f(n-1), f(0), \dots, f(n+j-1))$. Hence, the orbit $C_n(f)$ consists of all cyclic shifts of f .



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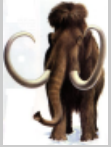
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Using the natural order $0 < 1$, the set $\{0, 1\}^n$ is totally ordered by the lexicographical order. For $f, g \in \{0, 1\}^n$ we say

$$f < g : \Leftrightarrow \exists i \in \underline{n} : f(j) = g(j) \text{ for } j < i \text{ and } f(i) < g(i).$$

We choose the smallest element of an orbit as its canonical representative.



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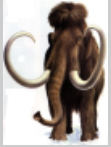
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0,1-Vectors (cont.)

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The stabilizer of $f \in \{0, 1\}^n$ is a subgroup of C_n , whence again a cyclic group. We call f **acyclic** if its stabilizer consists of the identity only. If f is acyclic, then the canonical representative of $C_n(f)$ is called a **Lyndon word**.

The function of the last example is acyclic, and (00110101) is a Lyndon word.



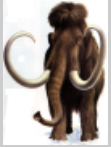
0,1-Vectors (cont.)

For example the orbit of $f = (10011010)$ under C_8 contains the vectors (10011010) , (00110101) , (01101010) , (11010100) , (10101001) , (01010011) , (10100110) , and (01001101) .

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We identify a 0,1-vector f with the set $f^{-1}(\{1\})$ of pre-images of 1. Then f is the characteristic function of $f^{-1}(\{1\})$. In this case we also assume that $f^{-1}(\{1\}) \subseteq \mathbb{Z}_n$.



Canons

A **canon** of length n consisting of $t \geq 1$ voices V_i is a set $\{V_1, \dots, V_t\}$ of 0,1-vectors $V_i \neq 0$ of length n , such that

1. $V_i \in C_n(V_1)$ for $1 \leq i \leq t$,
2. V_1 is acyclic,
3. the set of differences in $K := \bigcup_{i=1}^t V_i$ generates Z_n , i.e.

$$\langle K - K \rangle := \langle k - l \mid k, l \in K \rangle = Z_n.$$



Canons

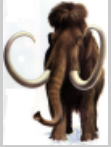
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Two canons $\{V_1, \dots, V_t\}$ and $\{W_1, \dots, W_s\}$ are called **isomorphic** if $s = t$ and if there exists some $\sigma \in C_n$ and a permutation τ in the symmetric group S_t such that $\sigma(V_i) = W_{\tau(i)}$ for $1 \leq i \leq t$.

Inner and Outer Rhythm of a Canon



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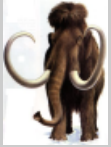
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The canon $\{V_1, \dots, V_t\}$ can be described as a pair (V_1, f) , where V_1 is the ***inner*** and f the ***outer rhythm*** of the canon. The inner rhythm describes the rhythm of any voice. The outer rhythm determines how the different voices are distributed over the n beats of a canon.

Inner and Outer Rhythm of a Canon



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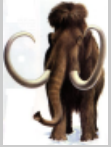
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For example consider $V_1 = (10011010)$, $V_2 = (01010011)$, and $V_3 = (11010100)$. We get a score of the form

$$\begin{array}{r} 10011010 \\ 01010011 \\ 11010100 \end{array}$$

Hence the outer rhythm of this canon is $f = (10010100)$.

Inner and Outer Rhythm of a Canon (cont.)



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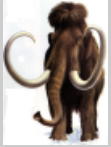
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Given a canon (V_1, f) of length n , there exists an isomorphic canon (L, f') where L is a Lyndon word, the canonical representative of $C_n(V_1)$, and f' is the canonical representative of $C_n(f)$.

Inner and Outer Rhythm of a Canon (cont.)



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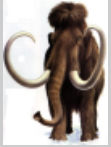
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Inner and Outer Rhythm of a Canon (cont.)



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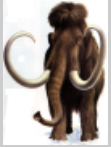
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Conversely, not each pair (L, f) where L is a Lyndon word and f a $0,1$ -vector determines a canon of length n .

Lemma. *Let $L \neq 0$ be a Lyndon word, and let f be a $0,1$ -vector both of length n . The pair (L, f) does not describe a canon if and only if there exists an integer $d > 1$ such that $d \mid n$, $d \mid k - l$ for all $k, l \in L$, and $d \mid k - l$ for all $k, l \in f$.*



Enumeration of Canons

Theorem. *The number of isomorphism classes of canons of length n is*

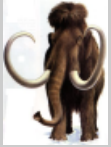
$$K_n := \sum_{d|n} \mu(d) \lambda(n/d) \alpha(n/d),$$

where μ is the **Moebius function**, $\lambda(1) = 1$,

$$\lambda(r) = \frac{1}{r} \sum_{s|r} \mu(s) 2^{r/s} \text{ for } r > 1,$$

$$\alpha(r) = \frac{1}{r} \sum_{s|r} \varphi(s) 2^{r/s} - 1 \text{ for } r \geq 1,$$

where φ is the **Euler totient function**.



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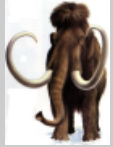
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The ideas of this proof and of the previous Lemma can be used in order to construct all rhythmical canons of length n .



Rhythmic Tiling Canons

In our last example of a canon, $((00110101), (01001010))$

10011010

01010011

11010100

we saw that there are beats where no, one, two, or three voices are playing at the same time.

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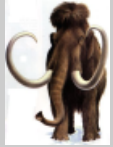
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A canon is called a **rhythmic tiling canon** if $Z_n = \bigcup_{i=1}^t V_i$. In other words, the voices are pairwise disjoint and cover entirely Z_n . The canon (L, f) is a tiling canon if and only if $L + f = Z_n$ and $|L| |f| = n$, thus Z_n is the direct sum of L and f .

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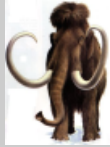
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For example $((00000101), (00110011))$ is the canon

01000001

10100000

00010100

00001010

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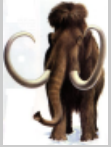
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Enumeration of Tiling Canons



Classification of rhythmic tiling canons by computing complete lists of their representatives. We have such lists for $n \leq 40$.

Comparison of T_n the number of nonisomorphic tiling canons and K_n the numbers of canons of length n .

n	T_n	K_n
2	1	1
3	1	5
4	2	13
5	1	41
6	3	110
7	1	341
8	6	1035
9	4	3298

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Enumeration of Tiling Canons

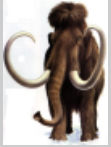
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n	T_n	K_n	n	T_n	K_n
2	1	1	10	6	10550
3	1	5	11	1	34781
4	2	13	12	23	117455
5	1	41	13	1	397529
6	3	110	14	13	1.370798
7	1	341	15	25	4.780715
8	6	1035	16	49	16788150
9	4	3298	17	1	59451809



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$n = 40$: $T_n = 64989$, $K_n = 755.578639.350274.265100$.

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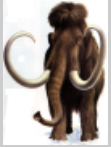
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Regular Complementary Canons of Maximal Category

A rhythmic tiling canon described by (L, f) is a ***regular complementary canon of maximal category*** (RCMC-canon) if both L and f are acyclic.

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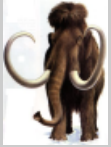
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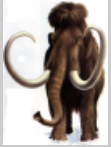
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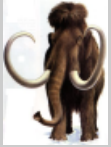


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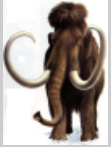
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Z_n is not a Hajós group if and only if n can be expressed in the form $p_1 p_2 n_1 n_2 n_3$ with p_1, p_2 primes, $n_i \geq 2$ for $1 \leq i \leq 3$, and $\gcd(n_1 p_1, n_2 p_2) = 1$.

How does such a canon sound?

[14, 8, 1, 5, 4, 4, 9, 9, 4, 6, 4, 9, 9, 4, 4, 5, 1, 8], [57, 12, 12, 3, 12, 12]



Vuza's Algorithm

If Z_n is not a Hajós group, Vuza presents an algorithm for constructing two acyclic vectors L and f of length n , such that $|L| = n_1n_2$, $|f| = p_1p_2n_3$, and $L + f = Z_n$.

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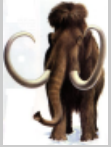
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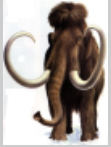
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Hence, both L or f can serve as the inner rhythm and the other one as the outer rhythm of an RCMC-canon.



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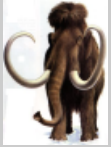
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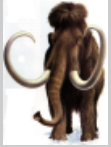
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Moreover, it is important to mention that there is some freedom for constructing these two sets, and each of these two sets can be constructed independently from the other one.

He also proves that when (L, f) describes an RCMC-canon, then also (kL, f) , (kL, kf) have this property for all $k \in Z_n^*$. (Here $kL = \{k\ell \mid \ell \in L\}$.)

Research by: G. Hajós, L. Rédei, N.G. de Bruijn, A.D. Sands.

Enumeration of Vuza Constructible Canons



An RCMC-canon which can be constructed by Vuza's algorithm and the previous remarks is called ***Vuza constructible canon***.

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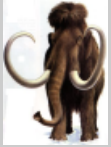
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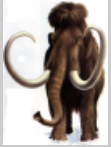
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An RCMC-canon which can be constructed by Vuza's algorithm and the previous remarks is called ***Vuza constructible canon***.

Enumeration of nonisomorphic Vuza constructible canons by construction:
This means that complete lists of these canons are available!



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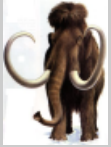
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n	p_1	p_2	n_1	n_2	n_3	$\#L$	$\#f$	$\#(L, f)$
72	2	3	2	3	2	3	6	18
108	2	3	2	3	3	3	180	540
120	2	3	2	5	2	16	20	320
120	2	5	2	3	2	8	6	48
144	2	3	4	3	2	6	36	216
144	2	3	2	3	4	3	2808	8424
168	2	3	2	7	2	104	42	4368
168	2	7	2	3	2	16	6	96
180	2	3	2	3	5	3	45360	136080
180	2	3	2	5	3	16	1000	16000
180	2	5	2	3	3	8	252	2016
180	2	3	5	3	2	9	60	540
180	3	5	3	2	2	6	12	72
200	2	5	2	5	2	125	20	2500
240	2	3	4	5	2	32	120	3840
240	2	5	4	3	2	16	36	576
240	2	3	2	5	4	16	26000	416000
240	2	5	2	3	4	8	6264	50112



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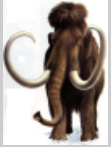
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Do there exist RCMC-canons which are not Vuza constructible?



Do there exist RCMC-canons which are not Vuza constructible?

The answer is:

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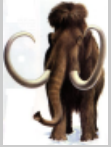
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Do there exist RCMC-canons which are not Vuza constructible?

The answer is: **YES!**

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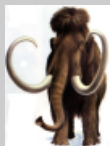
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Do there exist RCMC-canons which are not Vuza constructible?

The answer is: **YES!** For a long time it was an open question. I gave this answer at the Third International Seminar on Mathematical Music Theory and Music Informatics in 2002.

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Do there exist RCMC-canons which are not Vuza constructible?

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Idea: If you have an RCMC-canon, then subdivide each beat into 2 beats. Replace each quarter note or rest by 2 eighth notes or rests in the inner rhythm and enlarge the outer rhythm in a suitable way.

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Do there exist RCMC-canons which are not Vuza constructible?

The answer is: **YES!** For a long time it was an open question. I gave this answer at the Third International Seminar on Mathematical Music Theory and Music Informatics in 2002.

Idea: If you have an RCMC-canon, then subdivide each beat into 2 beats. Replace each quarter note or rest by 2 eighth notes or rests in the inner rhythm and enlarge the outer rhythm in a suitable way.

Construction: Let (L, f) be an RCMC-canon. Construct L' by replacing in L each occurrence of 1 by 11 and 0 by 00. And construct f' by replacing each 1 in f by 01 and 0 in f by 00. This way we construct from the RCMC-canon (L, f) of length n an RCMC-canon (L', f') of length $2n$.

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Among the 216 RCMC-canons of length $2 \cdot 72 = 144$ with $|L| = 12$ we did not find a canon which was constructed in this way from the 18 canons of length 72.

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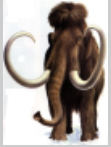
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We can do much better

Define a composition of two 0,1-vectors $f \in \{0, 1\}^{\underline{n}}$, $g \in \{0, 1\}^{\underline{m}}$. Then $f(g)$ is the vector in $\{0, 1\}^{\underline{nm}}$ where each 1 in f is replaced by g and each 0 in f is replaced by $0^m = 0 \dots 0$. Each $i \in \underline{nm}$ can uniquely be written as $i = qm + r$ with $q \in \underline{n}$ and $r \in \underline{m}$. Then

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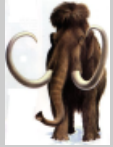
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Properties:

1. $\text{wt}(f(g)) = \text{wt}(f) \text{wt}(g)$.

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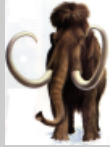
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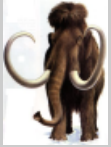
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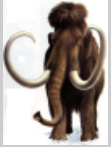
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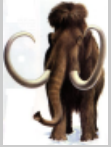
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6. If f and g are the canonical representatives of $C_n(f)$ and $C_m(g)$, then $f(g)$ is the canonical representative of $C_{nm}(f(g))$.



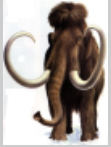
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6. If f and g are the canonical representatives of $C_n(f)$ and $C_m(g)$, then $f(g)$ is the canonical representative of $C_{nm}(f(g))$.
7. If f is a Lyndon word and g is the canonical representative of $C_m(g)$, then $f(g)$ is a Lyndon word.



New Canons from Old Canons

Theorem.

1. If (L_1, f_1) is a canon of length n_1 and (L_2, f_2) is a canon of length n_2 , then $(L_1(L_2), f_1(f_2))$ is a canon of length n_1n_2 .

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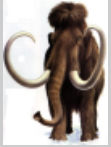
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New Canons from Old Canons



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2. If (L_1, f_1) is a tiling canon and (L_2, f_2) is a tiling canon, then $(L_1(L_2), f_1(f_2))$ is a tiling canon.

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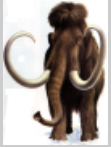
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3. If (L_1, f_1) is an RCMC-canon and (L_2, f_2) is a tiling canon, then $(L_1(L_2), f_1(f_2))$ is an RCMC-canon.

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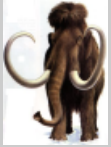
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Are these all RCMC-Canons?



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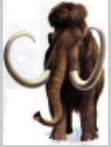
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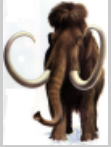
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The answer is:

Are these all RCMC-Canons?



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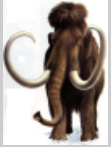
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The answer is: **NO!**



Are these all RCMC-Canons?

The answer is: **NO!**

Construction: Consider an RCMC-canon (L, f) of length n . For fixed L try to find all Lyndon words f' , such that $L + f' = Z_n$.

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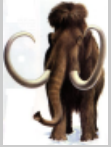
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Inspired by discussions with E. Amiot, I developed a backtracking algorithm for this problem. For given inner rhythm L it finds all outer rhythms f so that (L, f) is an RCMC-canon.

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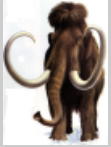
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Complete lists of RCMC-canons



If (L, f) is a tiling canon of length n , then the weights of L and f are divisors of n .

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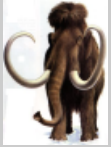
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Complete lists of RCMC-canons

If (L, f) is a tiling canon of length n , then the weights of L and f are divisors of n .

Theorem. (Sands) If $\text{wt}(L)$ or $\text{wt}(f)$ is a prime power, then f has cyclic symmetries. Thus (L, f) is not an RCMC-canon.

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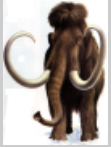
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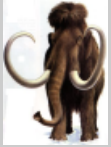
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Construction: Find all suitable decompositions of n into two integers, $n = rs$. If r and s are not powers of a prime assume that $r \leq s$ and continue with the following construction. Determine all Lyndon words of length n and weight r over $\{0, 1\}$. In order to decrease their number, determine the number of orbits of these Lyndon words under the action of the affine group. For each of these representatives L determine all acyclic outer rhythms f , such that $L + f = Z_n$.

All RCMC-canons of length 72 and 108



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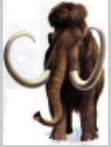
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We have to consider the two compositions $72 = 6 * 12$ and $108 = 6 * 18$.

There are 2.169882 Lyndon words of length 72 and weight 6 over $\{0, 1\}$.

There remain just 3 Lyndon words which can be extended to an RCMC-canon. For each Lyndon word there exist (the same 6) outer rhythms which can be used to determine an RCMC-canon. Thus all RCMC-canons of length 72 are Vuza constructible.

All RCMC-canons of length 72 and 108



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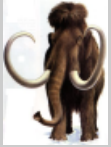
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There are 17.717859 Lyndon words of length 108 and weight 6 over $\{0, 1\}$. They are collected into 514754 orbits. There remains only one orbit representative which can be extended to an RCMC-canon. (This orbit contains 3 different Lyndon words.) For each of these Lyndon words there exist (the same 252) outer rhythms which can be used to determine an RCMC-canon. Thus some RCMC-canons of length 108 are not Vuza constructible.



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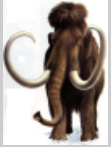
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Do there exist RCMC-canons which are not Vuza constructible?

We can do much better

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Are these all RCMC-Canons?

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